

# **AGENDA**

- Accelerated Computing Basics
- What are Compiler Directives?
- Accelerating Applications with OpenACC
  - Identifying Available Parallelism
  - Exposing Parallelism
  - Optimizing Data Locality
- Next Steps

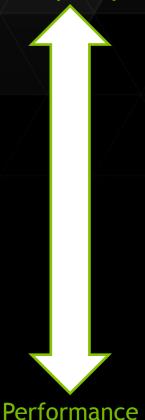
# ACCELERATED COMPUTING BASICS

# WHAT IS ACCELERATED COMPUTING?



# SIMPLICITY & PERFORMANCE

#### Simplicity



#### Accelerated Libraries

- Little or no code change for standard libraries; high performance
- Limited by what libraries are available

#### Compiler Directives

- High Level: Based on existing languages; simple and familiar
- ▶ High Level: Performance may not be optimal

#### Parallel Language Extensions

- Expose low-level details for maximum performance
- Often more difficult to learn and more time consuming to implement

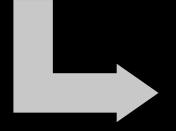
# CODE FOR PORTABILITY & PERFORMANCE

Libraries

• Implement as much as possible using portable libraries.

Directives

• Use directives to implement portable code.



Languages

 Use lower level languages for important kernels.

# WHAT ARE COMPILER DIRECTIVES?

# WHAT ARE COMPILER DIRECTIVES?

```
int main() {
  do_serial_stuff()

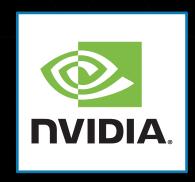
for(int i=0; i < BIGN; i++)
  {
    ...compute intensive work
  }

  do_more_serial_stuff();
}</pre>
```

Programmer inserts compiler hints.

Execution Begins on the CPU.

DataComopElececuteiner artes & Pld Clock & PU.



Data and Execution returns to the CPU.

# OPENACC: THE STANDARD FOR GPU DIRECTIVES

Simple: Easy path to accelerate compute intensive applications

Open: Open standard that can be implemented anywhere

Portable: Represents parallelism at a high level making it portable to any architecture



# OPENACC MEMBERS AND PARTNERS

































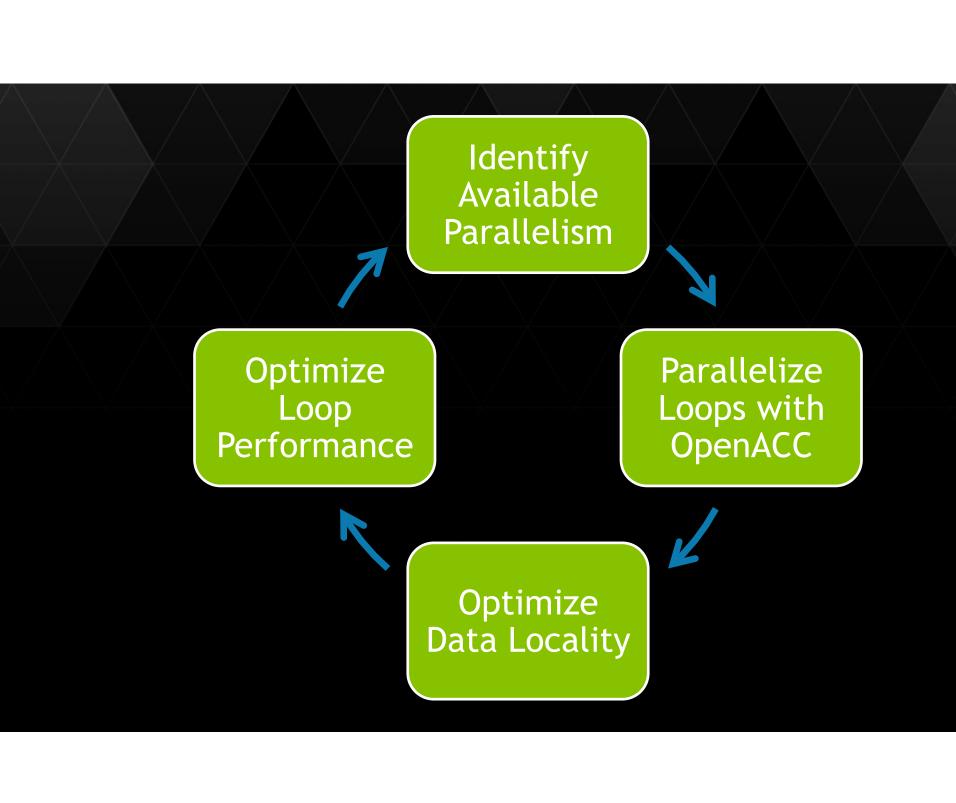








# ACCELERATING APPLICATIONS WITH OPENACC



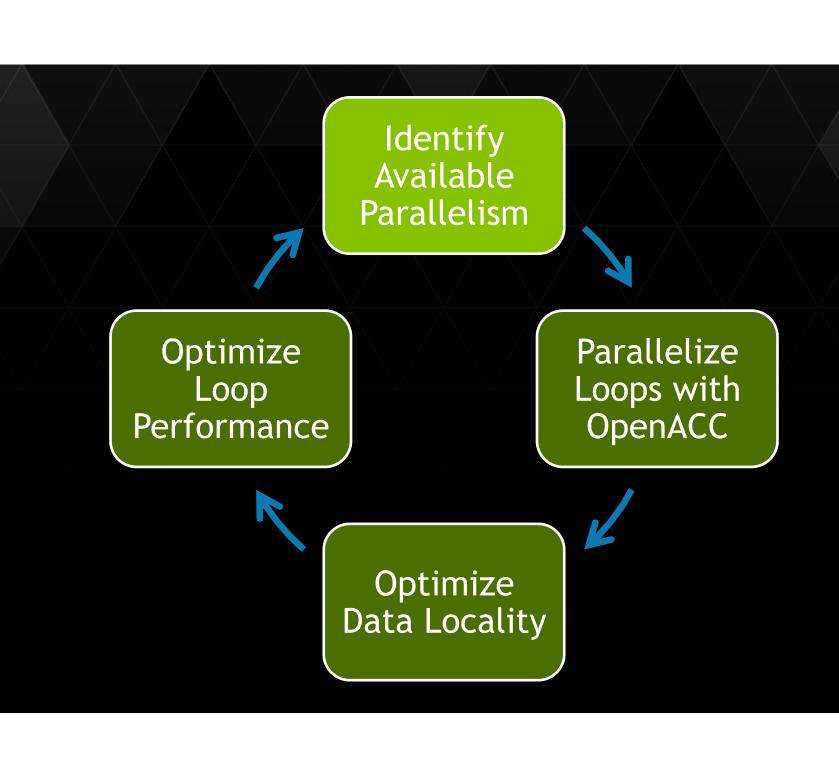
# **EXAMPLE: JACOBI ITERATION**

Iteratively converges to correct value (e.g. Temperature), by computing new values at each point from the average of neighboring points.

Common, useful algorithm A(i,j+1)Example: Solve Laplace equation in 2D:  $\nabla t = f(x,y) = 0$  A(i-1,j) A(i+1,j) A(i,j+1) A(i,j+1) A(i,j+1) A(i,j+1)

# JACOBI ITERATION: C CODE

```
while ( err > tol && iter < iter max ) {</pre>
                                                                       Iterate until converged
  err=0.0;
                                                                        Iterate across matrix
 for( int j = 1; j < n-1; j++) {
                                                                              elements
    for (int i = 1; i < m-1; i++) {
                                                                      Calculate new value from
      Anew[j][i] = 0.25 * (A[j][i+1] + A[j][i-1] +
                                                                             neighbors
                            A[j-1][i] + A[j+1][i]);
      err = max(err, abs(Anew[j][i] - A[j][i]));
                                                                       Compute max error for
                                                                            convergence
 for( int j = 1; j < n-1; j++) {
    for( int i = 1; i < m-1; i++ ) {
                                                                      Swap input/output array
      A[j][i] = Anew[j][i];
  iter++;
                                                                         14
```



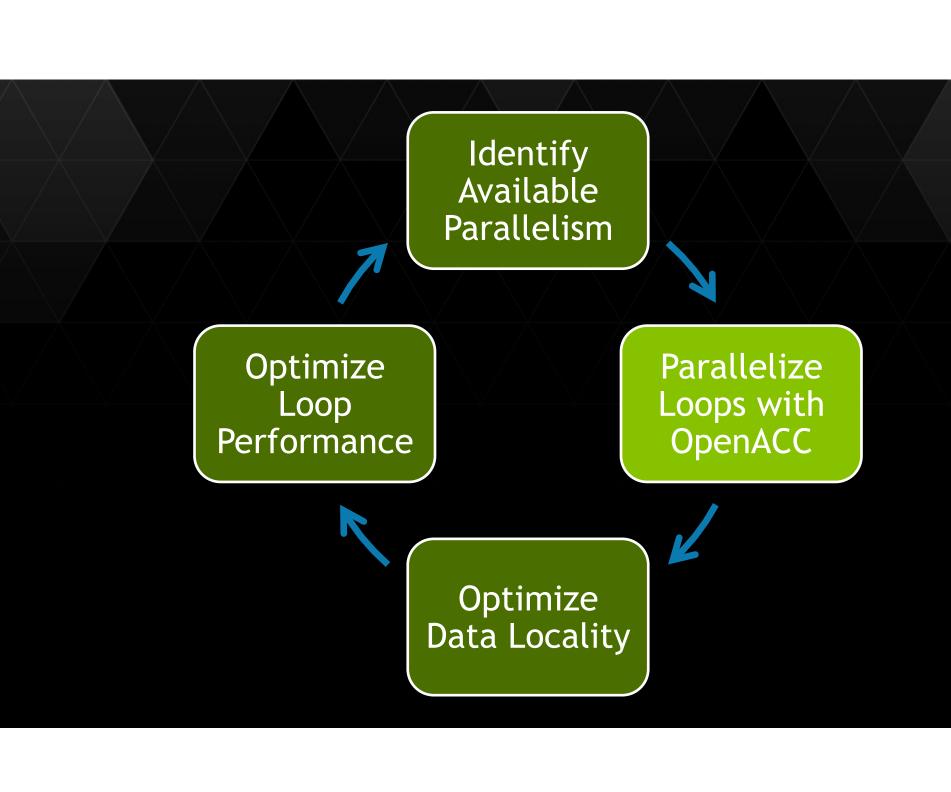
# IDENTIFY PARALLELISM

```
while ( err > tol && iter < iter max ) {
  err=0.0;
  for ( int j = 1; j < n-1; j++) {
    for (int i = 1; i < m-1; i++) {
      Anew[j][i] = 0.25 * (A[j][i+1] + A[j][i-1] +
                           A[j-1][i] + A[j+1][i]);
      err = max(err, abs(Anew[j][i] - A[j][i]));
  for( int j = 1; j < n-1; j++) {
    for ( int i = 1; i < m-1; i++ ) {
      A[j][i] = Anew[j][i];
  iter++;
```

Data dependency between iterations.

Independent loop iterations

Independent loop iterations



# OPENACC DIRECTIVE SYNTAX

```
► C/C++
```

```
#pragma acc directive [clause [,] clause] ...]
...often followed by a structured code block
```

Fortran

```
!$acc directive [clause [,] clause] ...]
often paired with a matching end directive surrounding a struct
```

...often paired with a matching end directive surrounding a structured code block:

!\$acc end directive



# OPENACC PARALLEL LOOP DIRECTIVE

parallel - Programmer identifies a block of code containing parallelism.
Compiler generates a kernel.

loop - Programmer identifies a loop that can be parallelized within the kernel.

NOTE: parallel & loop are often placed together

```
#pragma acc parallel loop
for(int i=0; i<N; i++)
{
    y[i] = a*x[i]+y[i];
}</pre>
```

#### Kernel:

A function that runs in parallel on the GPU

# PARALLELIZE WITH OPENACC

```
while ( err > tol && iter < iter max ) {
  err=0.0;
#pragma acc parallel loop reduction(max:err)
  for ( int j = 1; j < n-1; j++) {
    for(int i = 1; i < m-1; i++) {
      Anew[j][i] = 0.25 * (A[j][i+1] + A[j][i-1] +
                           A[j-1][i] + A[j+1][i]);
      err = max(err, abs(Anew[j][i] - A[j][i]);
#pragma acc parallel loop
  for ( int j = 1; j < n-1; j++) {
    for( int i = 1; i < m-1; i++ ) {
      A[j][i] = Anew[j][i];
  iter++;
```

Parallelize loop on accelerator

Parallelize loop on accelerator

\* A reduction means that all of the N\*M values for err will be reduced to just one, the max.

# BUILDING THE CODE

```
$ pgcc -fast -acc -ta=tesla -Minfo=all laplace2d.c
main:
     40, Loop not fused: function call before adjacent loop
         Generated vector sse code for the loop
     51, Loop not vectorized/parallelized: potential early exits
     55, Accelerator kernel generated
         55, Max reduction generated for error
         56, #pragma acc loop gang /* blockIdx.x */
         58, #pragma acc loop vector(256) /* threadIdx.x */
     55, Generating copyout (Anew[1:4094][1:4094])
         Generating copyin(A[:][:])
         Generating Tesla code
     58, Loop is parallelizable
     66, Accelerator kernel generated
         67, #pragma acc loop gang /* blockIdx.x */
         69, #pragma acc loop vector(256) /* threadIdx.x */
     66, Generating copyin (Anew[1:4094][1:4094])
         Generating copyout(A[1:4094][1:4094])
         Generating Tesla code
     69, Loop is parallelizable
```

# OPENACC KERNELS DIRECTIVE

The kernels construct expresses that a region may contain parallelism and the compiler determines what can safely be parallelized.

```
#pragma acc kernels

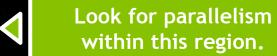
{
  for (int i=0; i<N; i++)
  {
    x[i] = 1.0;
    y[i] = 2.0;
  }

for (int i=0; i<N; i++)
  {
    y[i] = a*x[i] + y[i];
  }
  kernel 2</pre>
```

The compiler identifies 2 parallel loops and generates 2 kernels.

#### PARALLELIZE WITH OPENACC KERNELS

```
while ( err > tol && iter < iter max ) {</pre>
  err=0.0;
#pragma acc kernels
  for ( int j = 1; j < n-1; j++) {
    for(int i = 1; i < m-1; i++) {
      Anew[j][i] = 0.25 * (A[j][i+1] + A[j][i-1] +
                            A[j-1][i] + A[j+1][i]);
      err = max(err, abs(Anew[j][i] - A[j][i]));
  }
  for ( int j = 1; j < n-1; j++) {
    for( int i = 1; i < m-1; i++ ) {
      A[j][i] = Anew[j][i];
  iter++;
```



# BUILDING THE CODE

```
$ pgcc -fast -ta=tesla -Minfo=all laplace2d.c
main:
     40, Loop not fused: function call before adjacent loop
         Generated vector sse code for the loop
     51, Loop not vectorized/parallelized: potential early exits
     55, Generating copyout (Anew[1:4094][1:4094])
         Generating copyin(A[:][:])
         Generating copyout(A[1:4094][1:4094])
         Generating Tesla code
     57, Loop is parallelizable
     59, Loop is parallelizable
         Accelerator kernel generated
         57, #pragma acc loop gang /* blockIdx.y */
         59, #pragma acc loop gang, vector(128) /* blockIdx.x threadIdx.x */
         63, Max reduction generated for error
     67, Loop is parallelizable
     69, Loop is parallelizable
        Accelerator kernel generated
         67, #pragma acc loop gang /* blockIdx.y */
         69, #pragma acc loop gang, vector(128) /* blockIdx.x threadIdx.x */
```

# OPENACC PARALLEL LOOP VS. KERNELS

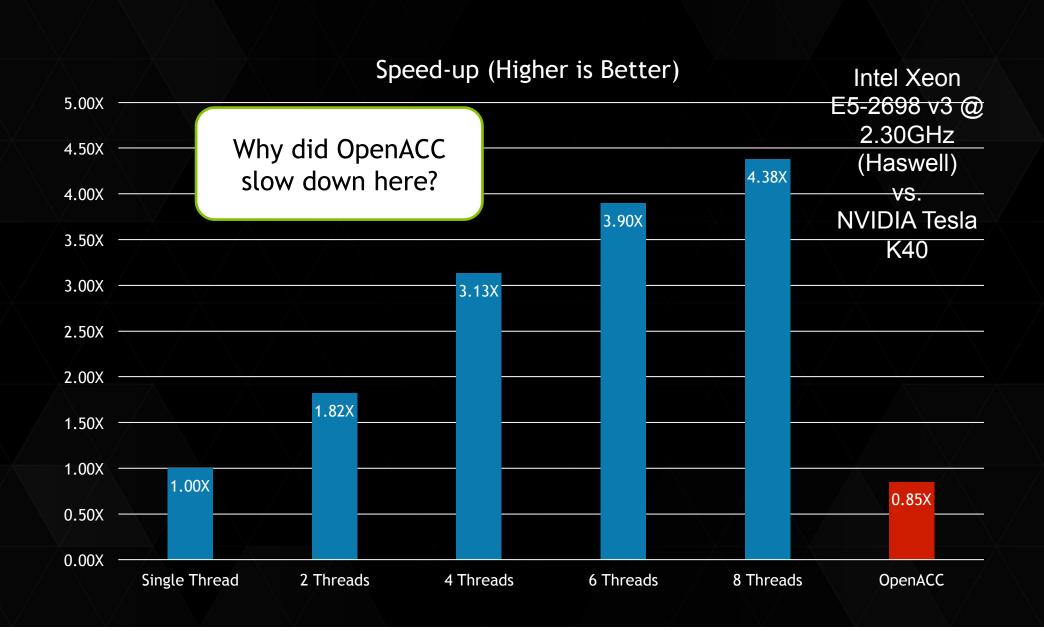
#### PARALLEL LOOP

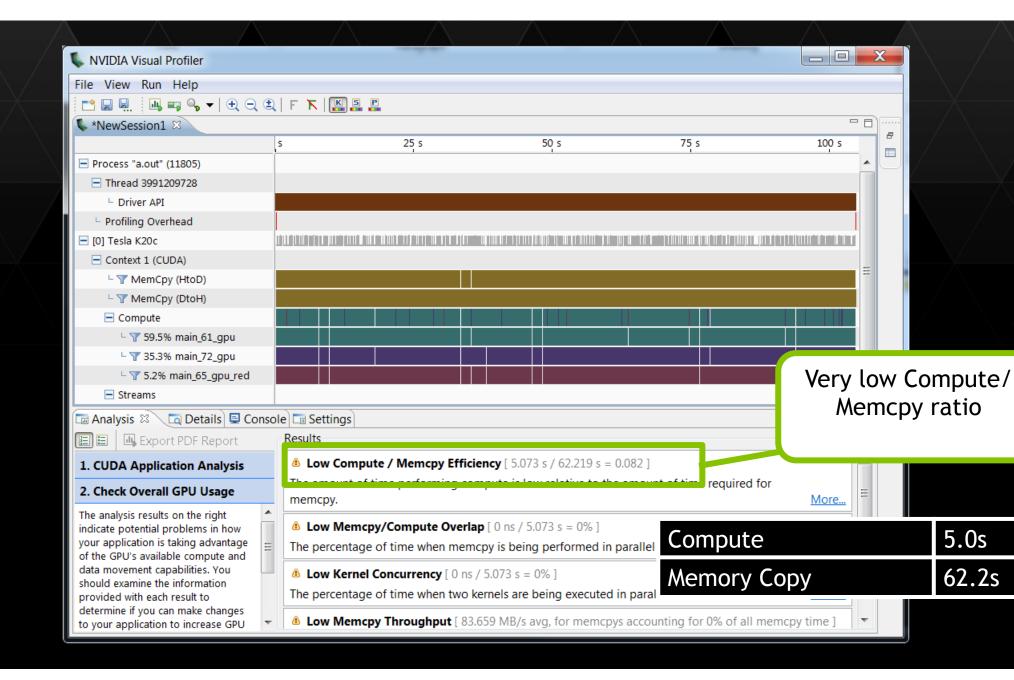
- Requires analysis by programmer to ensure safe parallelism
- Will parallelize what a compiler may miss
- Straightforward path from OpenMP

#### **KERNELS**

- Compiler performs parallel analysis and parallelizes what it believes safe
- Can cover larger area of code with single directive
- Gives compiler additional leeway to optimize.

Both approaches are equally valid and can perform equally well.





## EXCESSIVE DATA TRANSFERS

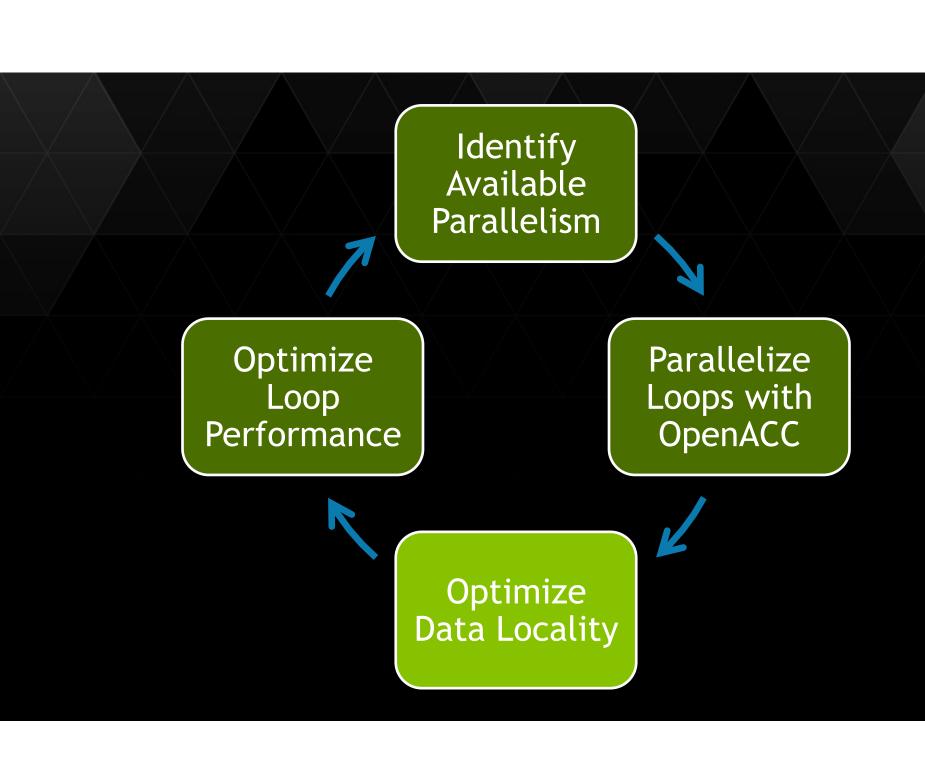
```
while ( err > tol && iter < iter max )</pre>
  err=0.0;
                                            #pragma acc kernels
            A. Anew resident
                                              A, Anew resident on
                 on host
                                                  accelerator
                                              for ( int j = 1; j < n-1; j++) {
                                                for (int i = 1; i < m-1; i++) {
            These copies
                                                  Anew[j][i] = 0.25 * (A[j][i+1] +
            happen every
                                                               A[j][i-1] + A[j-1][i] +
                                                               A[j+1][i]);
           iteration of the
                                                  err = max(err, abs(Anew[j][i] -
             outer while
                                                                      A[j][i]);
                 loop!
           A, Anew resident
                                                   A, Anew resident on
                on host
                                                       accelerator
```

### IDENTIFYING DATA LOCALITY

```
while ( err > tol && iter < iter max ) {</pre>
  err=0.0;
#pragma acc kernels
  for ( int j = 1; j < n-1; j++) {
    for(int i = 1; i < m-1; i++) {
      Anew[j][i] = 0.25 * (A[j][i+1] + A[j][i-1] +
                            A[j-1][i] + A[j+1][i]);
      err = max(err, abs(Anew[j][i] - A[j][i]));
  for ( int j = 1; j < n-1; j++) {
    for( int i = 1; i < m-1; i++ ) {
      A[j][i] = Anew[j][i];
  iter++;
```

Does the CPU need the data between these loop nests?

Does the CPU need the data between iterations of the convergence loop?



# DEFINING DATA REGIONS

The data construct defines a region of code in which GPU arrays remain on the GPU and are shared among all kernels in that region.

```
#pragma acc data
{
#pragma acc parallel loop
...
#pragma acc parallel loop
...
}
```

**Data Region** 

Arrays used within the data region will remain on the GPU until the end of the data region.

# DATA CLAUSES

```
Allocates memory on GPU and copies data from
copy ( list)
                host to GPU when entering region and copies
                data to the host when exiting region.
copyin ( list ) Allocates memory on GPU and copies data from
                host to GPU when entering region.
copyout (list) Allocates memory on GPU and copies data to
                the host when exiting region.
create ( list ) Allocates memory on GPU but does not copy.
present ( list ) Data is already present on GPU from another
                containing data region.
and present or copy[in|out], present or create, deviceptr.
```

# ARRAY SHAPING

- Compiler sometimes cannot determine size of arrays
  - Must specify explicitly using data clauses and array "shape"

```
<u>C/C++</u>
```

```
#pragma acc data copyin(a[0:size]) copyout(b[s/4:3*s/4])
Fortran
!$acc data copyin(a(1:end)) copyout(b(s/4:3*s/4))
```

▶ Note: data clauses can be used on data, parallel, or kernels

#### OPTIMIZE DATA LOCALITY

```
#pragma acc data copy(A) create(Anew)
while ( err > tol && iter < iter max ) {</pre>
  err=0.0;
#pragma acc kernels
  for ( int j = 1; j < n-1; j++) {
    for(int i = 1; i < m-1; i++) {
      Anew[j][i] = 0.25 * (A[j][i+1] + A[j][i-1] +
                           A[j-1][i] + A[j+1][i]);
      err = max(err, abs(Anew[j][i] - A[j][i]));
  for ( int j = 1; j < n-1; j++) {
    for( int i = 1; i < m-1; i++ ) {
      A[j][i] = Anew[j][i];
  iter++;
```

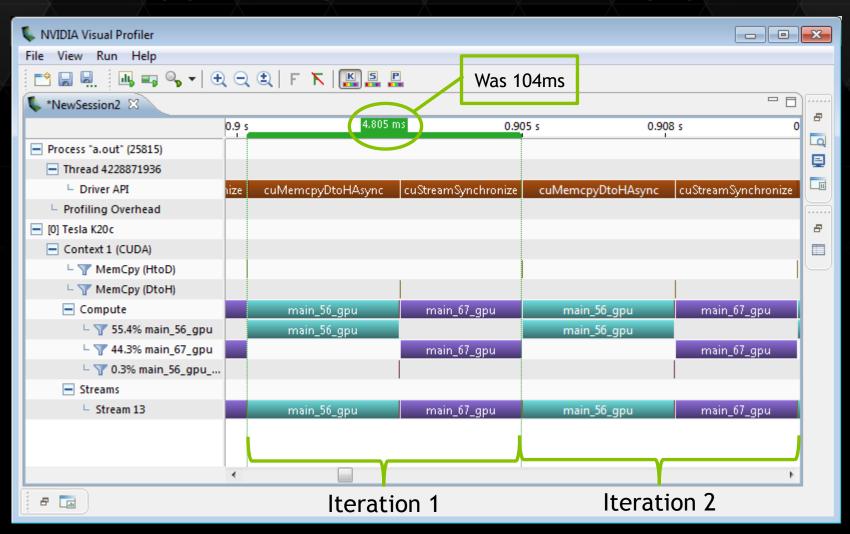
Copy A to/from the accelerator only when needed.

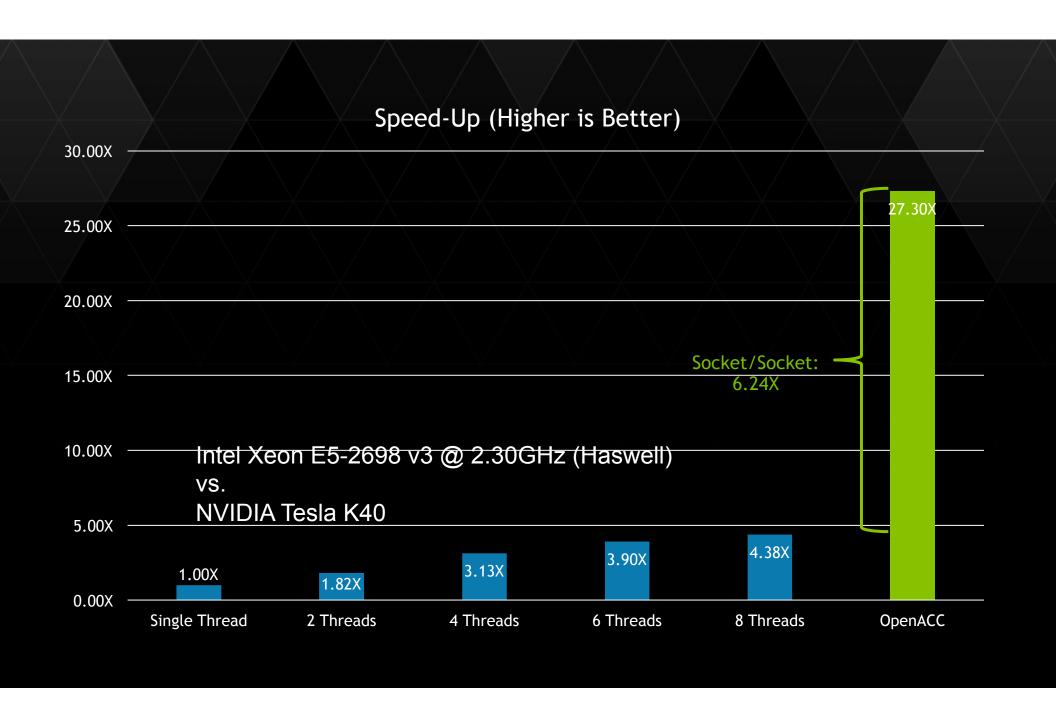
Create Anew as a device temporary.

# REBUILDING THE CODE

```
$ pgcc -fast -acc -ta=tesla -Minfo=all laplace2d.c
main:
     40, Loop not fused: function call before adjacent loop
         Generated vector sse code for the loop
     51, Generating copy(A[:][:])
         Generating create (Anew[:][:])
         Loop not vectorized/parallelized: potential early exits
     56, Accelerator kernel generated
         56, Max reduction generated for error
         57, #pragma acc loop gang /* blockIdx.x */
         59, #pragma acc loop vector(256) /* threadIdx.x */
     56, Generating Tesla code
     59, Loop is parallelizable
     67, Accelerator kernel generated
         68, #pragma acc loop gang /* blockIdx.x */
         70, #pragma acc loop vector(256) /* threadIdx.x */
     67, Generating Tesla code
     70, Loop is parallelizable
```

## VISUAL PROFILER: DATA REGION





#### OPENACC PRESENT CLAUSE

It's sometimes necessary for a data region to be in a different scope than the compute region.

When this occurs, the **present** clause can be used to tell the compiler data is already on the device.

Since the declaration of A is now in a higher scope, it's necessary to shape A in the present clause.

High-level data regions and the present clause are often critical to good performance.

```
function main(int argc, char **argv)
{
    #pragma acc data copy(A)
    {
        laplace2D(A,n,m);
    }
}
```

```
function laplace2D(double[N][M] A,n,m)
{
    #pragma acc data present(A[n][m]) create(Anew)
    while ( err > tol && iter < iter_max ) {
        err=0.0;
        ...
    }
}</pre>
```

Identify Available Parallelism

Optimize Loop Performance Parallelize Loops with OpenACC



Optimize

Data Locality



Watch S5195 - Advanced OpenACC Programming on gputechconf.com

# **NEXT STEPS**

#### 1. Identify Available Parallelism

What important parts of the code have available parallelism?

#### 2. Parallelize Loops

- Express as much parallelism as possible and ensure you still get correct results.
- Because the compiler *must* be cautious about data movement, the code will generally slow down.

#### 3. Optimize Data Locality

The programmer will *always* know better than the compiler what data movement is unnecessary.

#### 4. Optimize Loop Performance

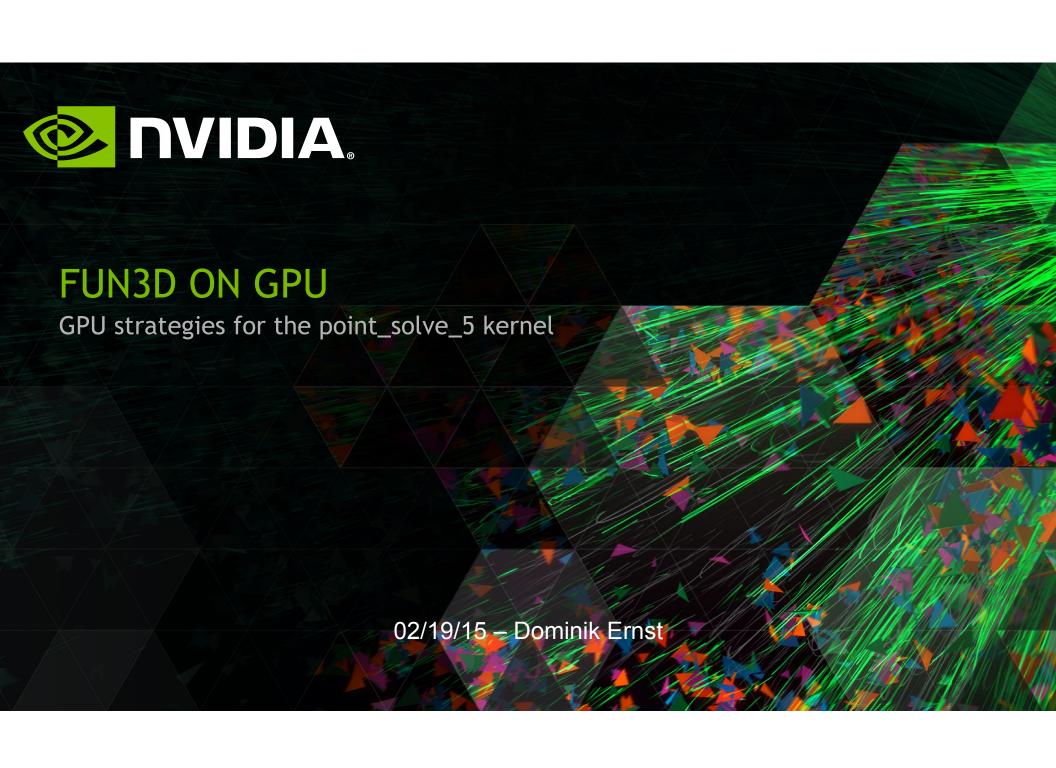
Don't try to optimize a kernel that runs in a few us or ms until you've eliminated the excess data motion that is taking many seconds.

Step 1 Step 2 Step 3 Step 4 **Identify Available Optimize Data** Parallelize Loops Optimize Loops **Parallelism** with OpenACC Locality **Application Speed-up Development Time** 

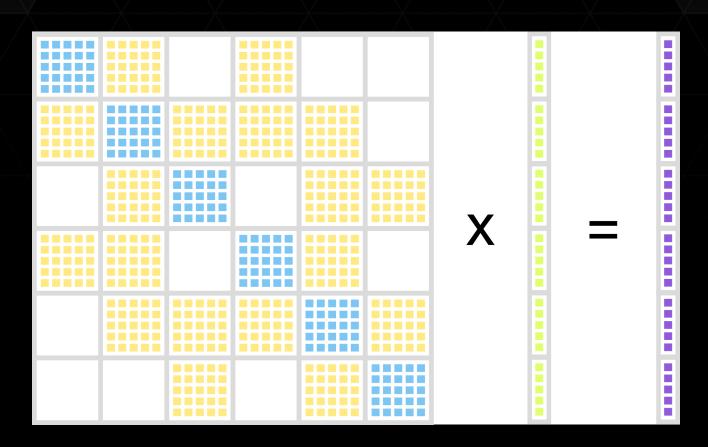
## FOR MORE INFORMATION

- Check out <a href="http://openacc.org/">http://openacc.org/</a>
- Watch tutorials at <a href="http://www.gputechconf.com/">http://www.gputechconf.com/</a>
- Share your successes at WACCPD at SC15.

Email me: jlarkin@nvidia.com



# POINT\_SOLVE\_5

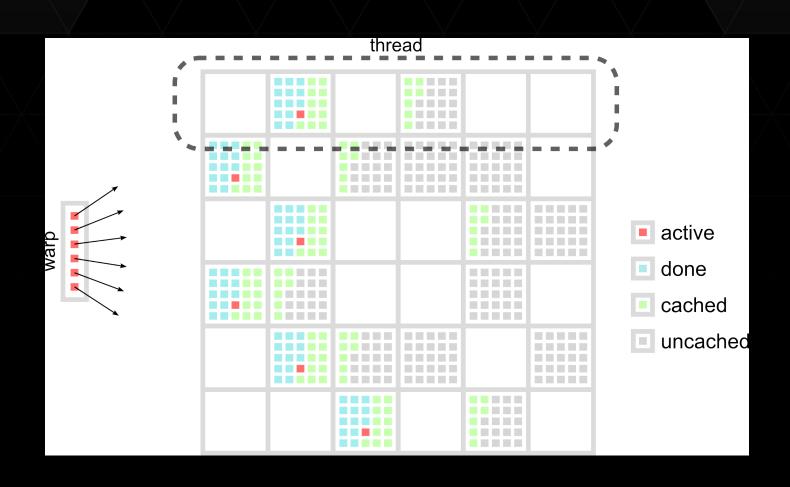


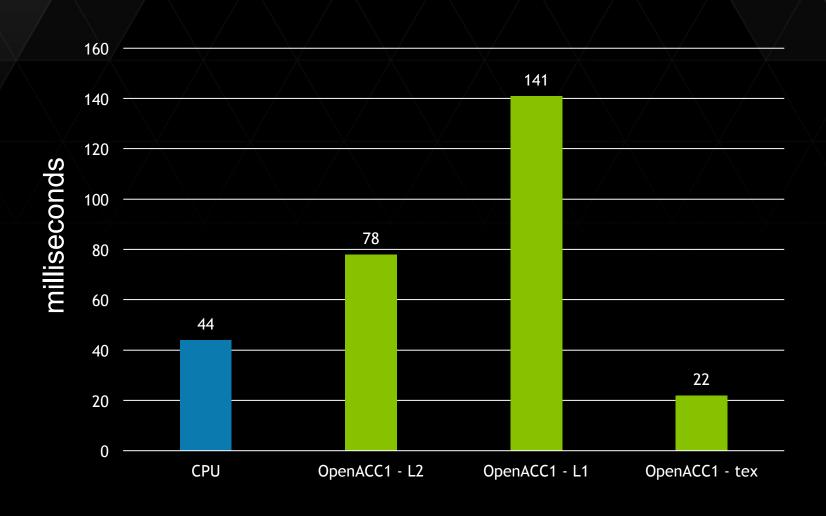
- ▶ CPU: One socket E5-2690 @ 3Ghz, 10 cores
- ▶ GPU: K40c, boost clocks, ECC off
- Dataset: DPW-Wing, 1M cells
- One call of point\_solve5 over all colors
- No transfers
- ▶ 1 CPU core: 300ms
- ▶ 10 CPU cores: 44ms

#### OPENACC1 - STRAIGHT FORWARD

```
!$acc parallel loop private(f1, f2, f3, f4, f5)
rhs solve : do n = start, end
    [...]
    istart = iam(n)
    iend = iam(n+1)
    do j = istart, iend
        icol = jam(j)
        f1 = f1 - a \circ ff(1,1,j) * dq(1,icol)
        f2 = f2 - a \circ ff(2,1,j)*dq(1,icol)
        [...22 lines]
         f5 = f5 - a \circ ff(5,5,j)*dq(5,icol)
    end do
    [...]
end do
```

# OPENACC1 - A\_OFF ACCESS PATTERN

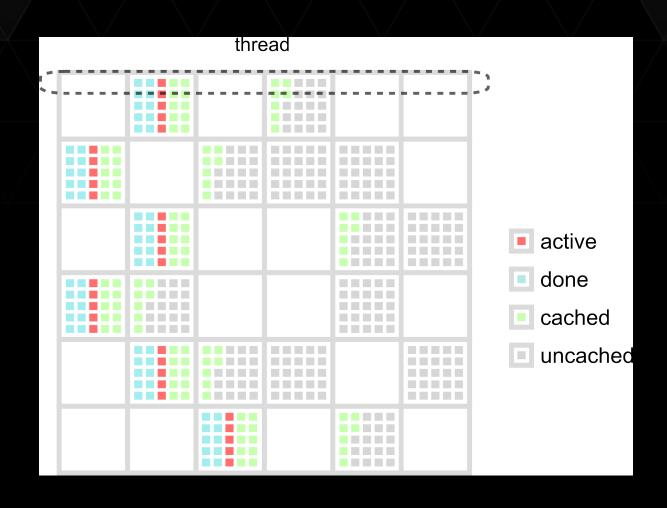


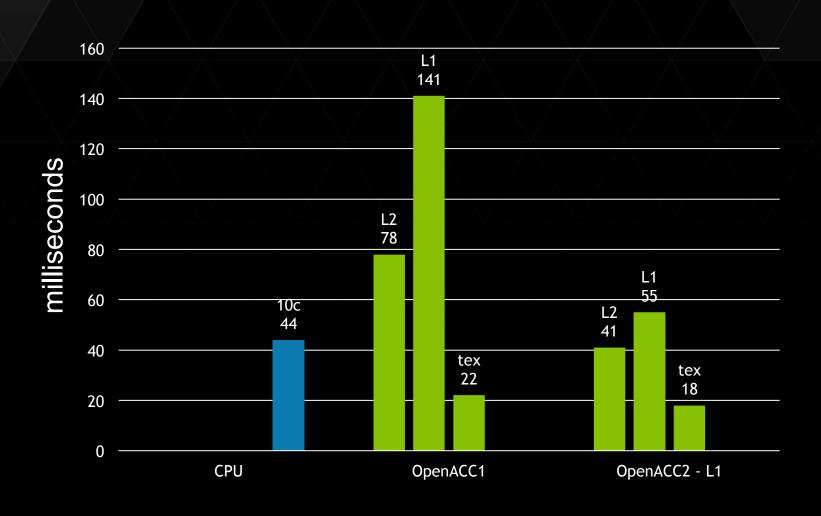


## OPENACC2 - REFORMULATED

```
!$acc parallel loop collapse(2) private(fk)
rhs solve : do n = start, end
    do k = 1,5
        T...1
        istart = iam(n)
        iend = iam(n+1)
        do j = istart, iend
            icol = jam(j)
            fk = fk - a off(k,1,j)*dq(1,icol)
            [... 3 lines]
            fk = fk - a \circ ff(k,5,j) *dq(1,icol)
        end do
    end do
    dq(k,n) = fk
end do
[Split off fw/bw substitution in extra loop]
```

# OPENACC2 - A\_OFF ACCESS PATTERN





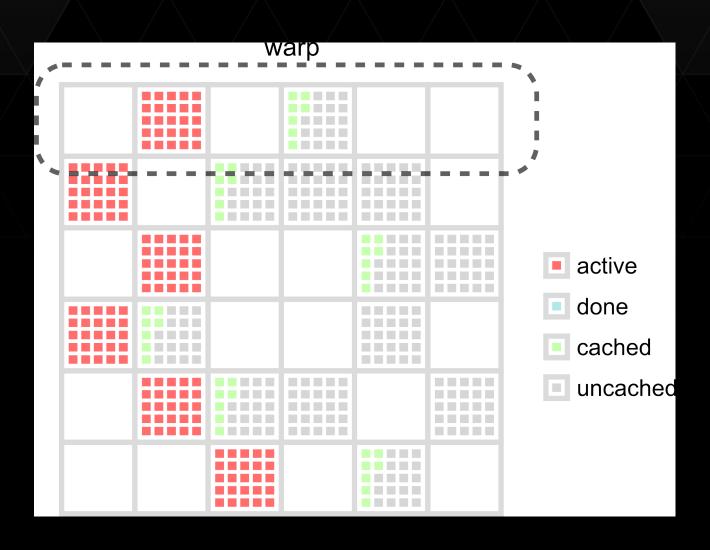
## CUDA FORTRAN - ADVANTAGES

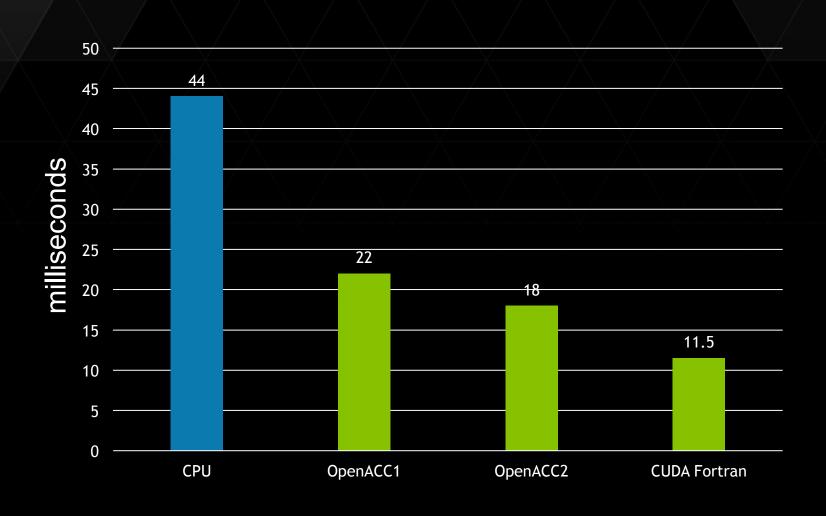
- Shared Memory: as explicitly managed cache and for cooperative reuse
- Inter thread communication in a thread block with shared memory
- Inter thread communication in a warp with \_\_shfl() instrinsic

## CUDA FORTRAN - 25 WIDE

```
! Calculate n, l, k based on threadIdx
! Loop over a off entries
istart = iam(n)
iend = iam(n+1)
do j = istart, iend
   ftemp = ftemp - a off(k,l,j)*dq(l,jam(j))
end do
! Reduction along the rows
fk = ftemp - shfl(ftemp, k+1*5)
fk = fk - shfl(ftemp, k+2*5)
fk = fk - shfl(ftemp, k+3*5)
fk = fk - shfl(ftemp, k+4*5)
```

# CUDA FORTRAN - 25 WIDE





## FUN3D CONCLUSIONS

- Unchanged code with OpenACC: 2.0x
- Modified code with OpenACC: 2.4x, modified code runs 50% slower on CPUs
- Highly optimized CUDA version: 3.7x
- Compiler options (e.g. how memory is accessed) have huge influence on OpenACC results
- Possible compromise: CUDA for few hotspots, OpenACC for the rest
- Very good OpenACC/CUDA interoperability: CUDA can use buffers managed with OpenACC data clauses
- Unsolved problem: data transfer in a partial port cause overhead